

# **A Comparison of SAT Encodings for Acyclicity of Directed Graphs**

Neng-Fa Zhou (CUNY & Relational AI)

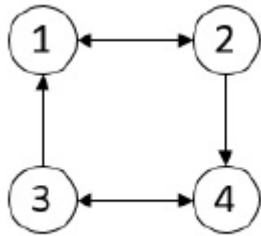
Ruiwei Wang (NUS)

Roland H.C. Yap (NUS)

## Synthesizing Acyclic Directed Graphs

- Given a base directed graph, find a subgraph that is acyclic and satisfies some side constraints.
  - Utility networks (water, electricity, gas, ...)
  - Planning problems
  - Markov networks
  - Neural networks
  - Bayesian networks

# The $\text{acyclic\_d}(V, E)$ Constraint



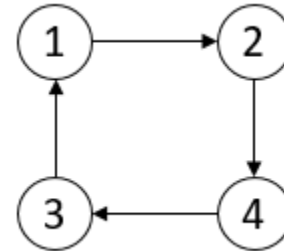
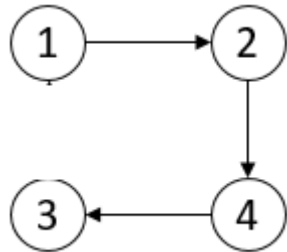
$\text{acyclic\_d}(V, E)$

$V = [v1, v2, v3, v4]$

$E = [(v1, v2), (v2, v1), (v2, v4),$   
 $(v3, v1), (v3, v4), (v4, v3)]$

$b(v) = 1$  iff  $v$  is in the subgraph

$b((v_i, v_j)) = 1$  iff  $(v_i, v_j)$  is in the subgraph



## SAT Encodings of the $\text{acyclic\_d}(V, E)$ Constraint

- Transitive closure encoding (TCE)
  - A graph is acyclic iff its transitive closure is acyclic.
  - $O(n^3)$  encoding size 😞
- Vertex-elimination encoding (VEE) [RankoohR22]



- If  $G'$  is cyclic, then  $G$  is cyclic
- Much more compact than TCE for sparse graphs 😊
- Asymptotically the same as TCE 😞

## SAT Encodings of the $\text{acyclic\_d}(V, E)$ Constraint (Cont.)

- Leaf-elimination encoding (LEE)
  - A graph is acyclic if the graph can be reduced to empty after leaves are repeatedly eliminated.
$$G_0 = G \xrightarrow{\quad} G_1 \quad \dots \quad \xrightarrow{\quad} G_t$$
  - Use a time variable for each vertex
  - The same as *tree-reduction encoding* [GebserJR14]
- LEE-u (use unary encoding for the time variable)
- LEE-b (use binary encoding for the time variable)
  - Compact 😊
  - Weak propagation 😞
  - The same as *binary labeling* [JanotaGM17]

## SAT Encodings of the $\text{acyclic\_d}(V, E)$ Constraint (Cont.)

- Hybrid encoding (HYB)
  - Start with VEE
  - Switch to LEE-b when the graph is dense
  - Enjoys VEE's strong propagation and LEE-b's conciseness 😊

## Vertex-Elimination Encoding (VEE)

- Vertex elimination operation

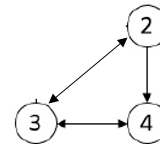
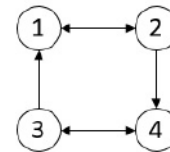
$$G_0 = ([v_1, v_2, \dots, v_n], E_0)$$

:

$$G_{i-1} = ([v_i, \dots, v_n], E_{i-1})$$



$$G_i = ([v_{i+1}, \dots, v_n], E_i)$$



- Encoding
  - $b((u, v_i)) \wedge b((v_i, w)) \rightarrow b((u, w))$
  - $\neg b((u, v_i)) \vee \neg b((v_i, u))$
- Encoding size:  $O(n^2)$ ,  $O(n \times d^2)$

## *Leaf-Elimination Encoding*

### *Using Unary Encoding for Time Variables (LEE-u)*

- Variables

- $A_{v,t}, v \in \{v_1, v_2, \dots, v_n\}, t \in 0..m$
- $A_{v,t} = 1$  iff vertex  $v$  has been eliminated by time  $t$

- Constraints

$$\text{For } v \in V: \sum_{u \in \text{nbr}^+(v)} b((v, u)) = 0 \leftrightarrow A_{v,0} \quad (1)$$

$$\text{For } v \in V, t \in 1..m: A_{v,t-1} \rightarrow A_{v,t} \quad (2)$$

$$\text{For } v \in V, t \in 1..m, u \in \text{nbr}^+(v): b((v, u)) \wedge \neg A_{u,t-1} \rightarrow \neg A_{v,t} \quad (3)$$

$$\text{For } v \in V: A_{v,m} \quad (4)$$

- Encoding size:  $O(n \times m), O(n \times m \times d)$



## *Leaf-Elimination Encoding*

### *Using Binary Encoding for Time Variables (LEE-b)*

- Variables

- $T_v' \in 0..m$

- Constraints

For  $v \in V$ :  $\sum_{u \in nbs^+(v)} b((v, u)) = 0 \leftrightarrow T_v = 0$  (5)

For  $v \in V, u \in nbs^+(v)$ :  $b((v, u)) \rightarrow T_v > T_u$  (6)

- Encoding size:  $O(n \times \log_2 m)$ ,  $O(n \times \log_2 m \times d)$

## *Hybrid Encoding (HYB)*

- Start with VEE
- Switch to LEE-b when the graph is dense
  - we show the correctness of HYB

$$|E_0 \cup E_1 \cdots \cup E_i| \geq \min(2.3 \times |E_0|, 30 \times |V_0|)$$

## Experimental Results

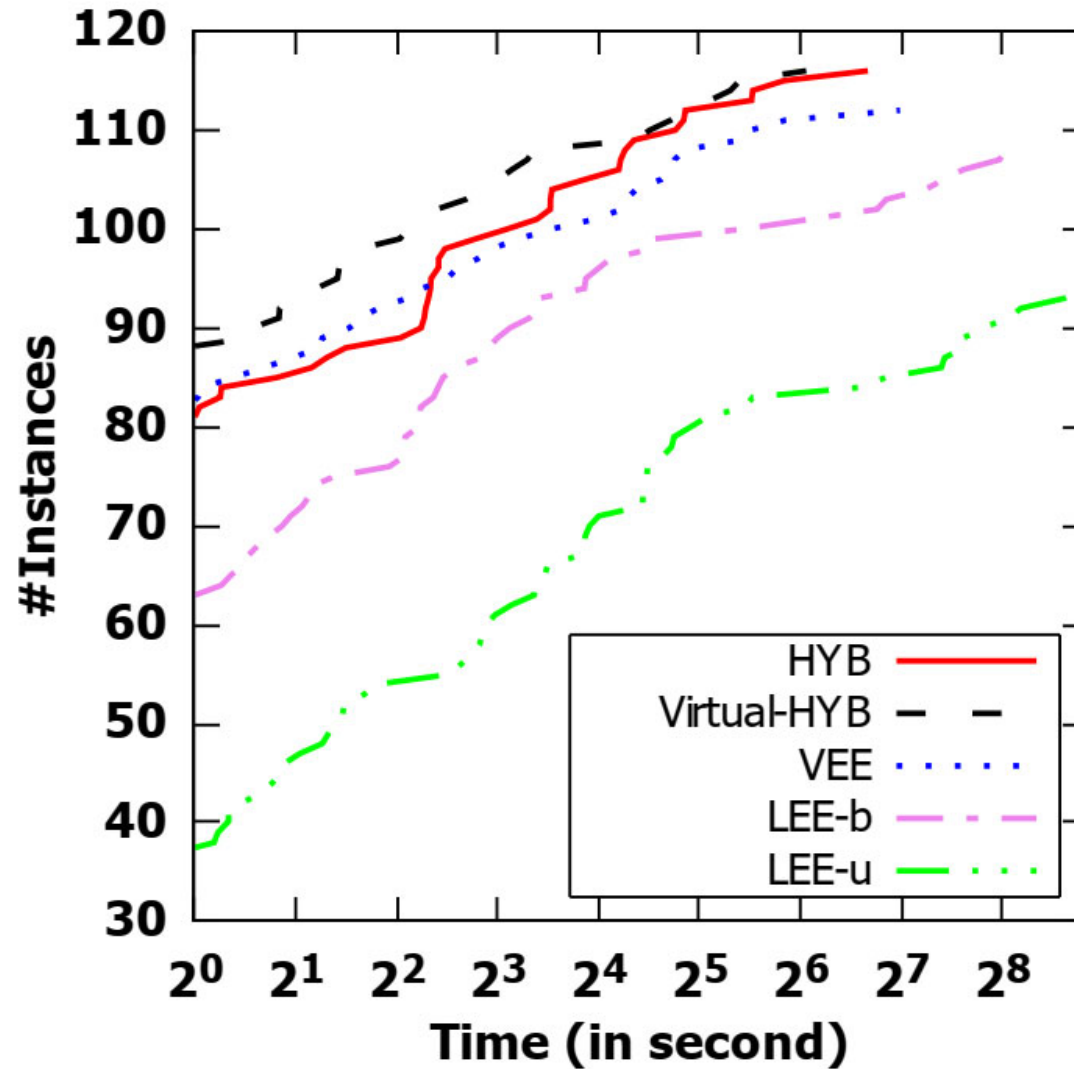
- Implementation
  - Picat ([picat-lang.org](http://picat-lang.org))
  - Kissat
- Benchmarks
  - GraphSAT benchmarks (n from 15 to 10002)
- Time limit: 10m per instance
- Platform
  - Ubuntu with a 3.20GHz and 16G RAM
  - Intel i7-8700 machine.

## Experimental Results (con.t)

#Insts	Benchmark	LEE-u	LEE-b	VEE	HYB	Virtual-HYB
26	COMB	28.03	42.54	<b>0.13</b>	<b>0.13</b>	<b>0.13</b>
31	EMPTYCORNER	160.89	1.99	3.84	1.99	<b>1.51</b>
36	EMPTYMIDDLE	67.01	6.93	0.92	2.02	<b>0.87</b>
9	ESCAPE	379.45	22.74	285.91	23.64	<b>15.38</b>
14	ROOMCHAIN	343.35	342.99	15.38	15.34	<b>11.94</b>
116	<b>TOTAL</b>	16350.64	6423.84	2944.2	565.29	<b>387.09</b>
	<b>AVE</b>	140.95	55.38	25.38	4.87	<b>3.34</b>
	<b>#SOLVED</b>	94	108	112	<b>116</b>	<b>116</b>

**Table 1** Summary of results

## Experimental Results (con.t)



(a) Runtime distribution

## *Conclusion*

- Contributions
  - A comparison of SAT encodings for acyclicity, including VEE, LEE-u, and LEE-b
  - Finding of a hybrid encoding, which combines VEE and LEE-b
- Future work
  - Heuristics for switching from VEE to LEE-b
  - Explore hybridization for encoding other constraints